

# Sequential Decision Aggregation: Accuracy and Decision Time

Sandra H. Dandach, Ruggero Carli and Francesco Bullo



Center for Control,  
Dynamical Systems & Computation  
University of California at Santa Barbara  
<http://motion.me.ucsb.edu>

MURI FA95500710528 Project Review: Behavioral Dynamics in  
Cooperative Control of Mixed Human/Robot Teams  
Center for Human and Robot Decision Dynamics, Aug 13, 2010

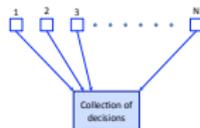
## Sequential decision aggregation: Outline

- 1 Setup & Literature Review
- 2 SDA: analysis of decision probabilities
- 3 SDA: scalability analysis of accuracy/decision time
- 4 Conclusions and future directions

## Setup & Literature Review

### Assumptions:

- 1  $N$  identical individuals, arbitrary local rule
- 2 Independent information
- 3 Aggregation of individual decisions



Group decision rule = SDA algorithm

$q$  out of  $N$  rule: decision as soon as  $q$  nodes report concordant opinion

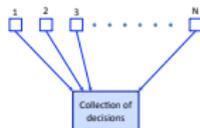
- Fastest rule fastest node decides for network ( $q = 1$ )
- Majority rule network agrees with majority decision ( $q = \lceil N/2 \rceil$ )

Goal #1: characterize decision probabilities of SDA  
as function of: threshold and SDM decision probabilities  
Goal #2: express accuracy & decision time  
as function of: decision threshold  $\times$  group size

## Setup & Literature Review

### Assumptions:

- 1  $N$  identical individuals, arbitrary local rule
- 2 Independent information
- 3 Aggregation of individual decisions



Group decision rule = SDA algorithm

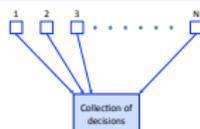
$q$  out of  $N$  rule: decision as soon as  $q$  nodes report concordant opinion

- Fastest rule fastest node decides for network ( $q = 1$ )
- Majority rule network agrees with majority decision ( $q = \lceil N/2 \rceil$ )

Goal #1: characterize decision probabilities of SDA  
as function of: threshold and SDM decision probabilities  
Goal #2: express accuracy & decision time  
as function of: decision threshold  $\times$  group size

## Assumptions:

- 1  $N$  identical individuals, arbitrary local rule
- 2 Independent information
- 3 Aggregation of individual decisions



## Group decision rule = SDA algorithm

$q$  out of  $N$  rule: decision as soon as  $q$  nodes report concordant opinion

- **Fastest rule** fastest node decides for network ( $q = 1$ )
- **Majority rule** network agrees with majority decision ( $q = \lceil N/2 \rceil$ )

- Goal #1: characterize decision probabilities of SDA  
as function of: threshold and SDM decision probabilities
- Goal #2: express accuracy & decision time  
as function of: decision threshold  $\times$  group size

## Distributed/decentralized detection

- 1 P. K. Varshney. *Distributed Detection and Data Fusion*. Signal Processing and Data Fusion. Springer Verlag, 1996
- 2 V. V. Veeravalli, T. Başar, and H. V. Poor. Decentralized sequential detection with sensors performing sequential tests. *Math Control, Signals & Systems*, 7(4):292–305, 1994
- 3 J. N. Tsitsiklis. Decentralized detection. In H. V. Poor and J. B. Thomas, editors, *Advances in Statistical Signal Processing*, volume 2, pages 297–344, 1993
- 4 J.-F. Chamberland and V. V. Veeravalli. Decentralized detection in sensor networks. *IEEE Trans Signal Processing*, 51(2):407–416, 2003

## Social networks

- 1 D. Acemoglu, M. A. Dahleh, I. Lobel, and A. Ozdaglar. Bayesian learning in social networks. Working Paper 14040, National Bureau of Economic Research, May 2008

## Literature review #2

For decentralized detection, with conditional independence of observations:

- Tsitsiklis '93: Bayesian decision problem with fusion center. For large networks identical local decision rules are asymptotically optimal
- Varshney '96: on non-identical decision rules with  $q$  out of  $N$ ,
  - 1 threshold rules are optimal at the nodes levels
  - 2 finding optimal thresholds requires solving  $N + 2^N$  equations
- Varshney '96: on optimal fusion rules for identical local decisions, " $q$  out of  $N$ " is optimal at the fusion center level

## Contributions today

- arbitrary decision makers (rather than optimal local rules)
- sequential aggregation (rather than "complete" aggregation)
- scalability analysis of accuracy / decision time

## Literature review #2

For decentralized detection, with conditional independence of observations:

- Tsitsiklis '93: Bayesian decision problem with fusion center. For large networks identical local decision rules are asymptotically optimal
- Varshney '96: on non-identical decision rules with  $q$  out of  $N$ ,
  - 1 threshold rules are optimal at the nodes levels
  - 2 finding optimal thresholds requires solving  $N + 2^N$  equations
- Varshney '96: on optimal fusion rules for identical local decisions, " $q$  out of  $N$ " is optimal at the fusion center level

## Contributions today

- arbitrary decision makers (rather than optimal local rules)
- sequential aggregation (rather than "complete" aggregation)
- scalability analysis of accuracy / decision time

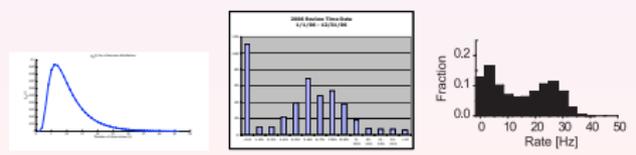
- 1 Setup & Literature Review
- 2 SDA: analysis of decision probabilities
- 3 SDA: scalability analysis of accuracy/decision time
- 4 Conclusions and future directions

Sequential decision maker (SDM)

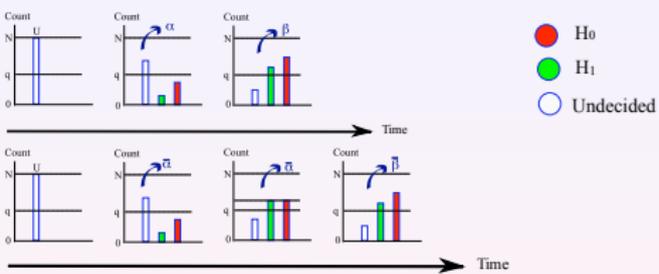
$$p_{ij}(t) := \text{Probability "say } H_i \text{ given } H_j \text{" at time } t$$

$$p_{ij} = \sum_{t=1}^{+\infty} p_{ij}(t), \quad E[T|H_i] = \sum_{t=1}^{+\infty} t(p_{1i}(t) + p_{0i}(t))$$

Assume knowledge of  $\{p_{ij}(t)\}_{t \in \mathbb{N}}$  for individual SDM, known exactly, calculated numerically, or measured empirically

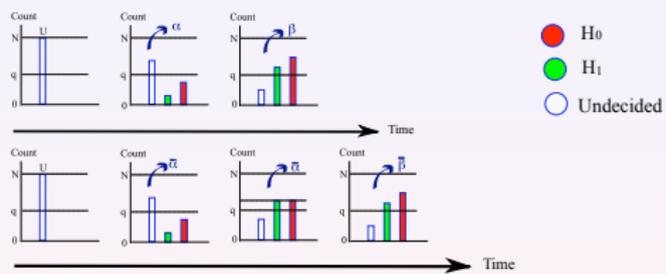


Sequential decision aggregation: Intermediate events



- aggregate states and divide in groups characterized by count
- calculate the probability of transition between the different groups
- characterize two states for network decisions  $H_0$  and  $H_1$

Sequential decision aggregation: Intermediate events



- aggregate states and divide in groups characterized by count
- calculate the probability of transition between the different groups
- characterize two states for network decisions  $H_0$  and  $H_1$

**Goal:** as function of SDM decision probabilities  $\{p_{ij}(t)\}_{t \in \mathbb{N}}$ ,  
compute SDA decision probabilities  $\{p_{ij}(t; N, q)\}_{t \in \mathbb{N}}$

General result:  $q$  out of  $N$  decision probabilities

$$p_{ij}(t; N, q) = \sum_{s_0=0}^{q-1} \sum_{s_1=0}^{q-1} \binom{N}{s_1 + s_0} \alpha(t-1, s_0, s_1) \beta_{ij}(t, s_0, s_1) + \sum_{s=q}^{\lfloor N/2 \rfloor} \binom{N}{2s} \bar{\alpha}(t-1, s) \bar{\beta}_{ij}(t, s)$$

As function of  $t$  and sizes, formulas for  $\alpha$ ,  $\beta$ ,  $\bar{\alpha}$ , and  $\bar{\beta}$   
computational complexity linear in  $N$

## Sequential decision aggregation: Computational approach

**Goal:** as function of SDM decision probabilities  $\{p_{ij}(t)\}_{t \in \mathbb{N}}$ ,  
compute SDA decision probabilities  $\{p_{ij}(t; N, q)\}_{t \in \mathbb{N}}$

General result:  $q$  out of  $N$  decision probabilities

$$p_{ij}(t; N, q) = \sum_{s_0=0}^{q-1} \sum_{s_1=0}^{q-1} \binom{N}{s_1 + s_0} \alpha(t-1, s_0, s_1) \beta_{ij}(t, s_0, s_1) + \sum_{s=q}^{\lfloor N/2 \rfloor} \binom{N}{2s} \bar{\alpha}(t-1, s) \bar{\beta}_{ij}(t, s)$$

As function of  $t$  and sizes, formulas for  $\alpha$ ,  $\beta$ ,  $\bar{\alpha}$ , and  $\bar{\beta}$   
computational complexity linear in  $N$

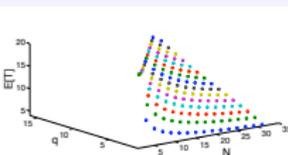
**Goal:** as function of SDM decision probabilities  $\{p_{ij}(t)\}_{t \in \mathbb{N}}$ ,  
compute SDA decision probabilities  $\{p_{ij}(t; N, q)\}_{t \in \mathbb{N}}$

General result:  $q$  out of  $N$  decision probabilities

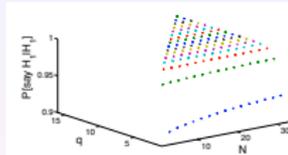
$$p_{ij}(t; N, q) = \sum_{s_0=0}^{q-1} \sum_{s_1=0}^{q-1} \binom{N}{s_1 + s_0} \alpha(t-1, s_0, s_1) \beta_{ij}(t, s_0, s_1) + \sum_{s=q}^{\lfloor N/2 \rfloor} \binom{N}{2s} \bar{\alpha}(t-1, s) \bar{\beta}_{ij}(t, s)$$

As function of  $t$  and sizes, formulas for  $\alpha$ ,  $\beta$ ,  $\bar{\alpha}$ , and  $\bar{\beta}$   
computational complexity linear in  $N$

## Illustration of results



Expected Decision Time



Probability of correct decision

- ( $H_0 : \mu = 0$ ) and ( $H_1 : \mu = 1$ )
- SPRT with  $p_f = p_m = 0.1$
- Gaussian noise  $\mathcal{N}(\mu, \sigma)$ ,  $\sigma = 1$  and  $\mu \in \{0, 1\}$

- 1 Setup & Literature Review
- 2 SDA: analysis of decision probabilities
- 3 SDA: scalability analysis of accuracy/decision time
- 4 Conclusions and future directions

Expected Decision Time:

$$\lim_{N \rightarrow \infty} \mathbb{E}[T|H_1, N, \text{fastest}] = \text{earliest possible decision time}$$

$$=: t_{min} = \min\{t \in \mathbb{N} \mid \text{either } p_{1|1}(t) \neq 0 \text{ or } p_{0|1}(t) \neq 0\}$$

Accuracy:

$$\lim_{N \rightarrow \infty} p_{0|1}(N, \text{fastest}) = \begin{cases} 0, & \text{if } p_{1|1}(t_{min}) > p_{0|1}(t_{min}) \\ 1, & \text{if } p_{1|1}(t_{min}) < p_{0|1}(t_{min}) \end{cases}$$

- 1 SDA accuracy is function of (SDM probability at  $t_{min}$ ), not of (SDA cumulative probability)!
- 2 hence, SDA accuracy is not monotonic with  $N$
- 3 hence, SDA accuracy is unrelated to SDM accuracy for large  $N$

Asymptotic results for the Fastest rule

Asymptotic results for the Majority rule

Expected Decision Time:

$$\lim_{N \rightarrow \infty} \mathbb{E}[T|H_1, N, \text{fastest}] = \text{earliest possible decision time}$$

$$=: t_{min} = \min\{t \in \mathbb{N} \mid \text{either } p_{1|1}(t) \neq 0 \text{ or } p_{0|1}(t) \neq 0\}$$

Accuracy:

$$\lim_{N \rightarrow \infty} p_{0|1}(N, \text{fastest}) = \begin{cases} 0, & \text{if } p_{1|1}(t_{min}) > p_{0|1}(t_{min}) \\ 1, & \text{if } p_{1|1}(t_{min}) < p_{0|1}(t_{min}) \end{cases}$$

- 1 SDA accuracy is function of (SDM probability at  $t_{min}$ ), not of (SDA cumulative probability)!
- 2 hence, SDA accuracy is not monotonic with  $N$
- 3 hence, SDA accuracy is unrelated to SDM accuracy for large  $N$

Expected Decision Time: Assume  $p_{1|1} > p_{0|1}$  and define

$$t_{< \frac{1}{2}} := \max\{t \in \mathbb{N} \mid p_{1|1}(0) + \dots + p_{1|1}(t) < 1/2\},$$

$$t_{> \frac{1}{2}} := \min\{t \in \mathbb{N} \mid p_{1|1}(0) + \dots + p_{1|1}(t) > 1/2\}$$

Then

$$\lim_{N \rightarrow \infty} \mathbb{E}[T|H_1, N, \text{majority}] = \frac{1}{2}(t_{< \frac{1}{2}} + t_{> \frac{1}{2}} + 1)$$

Accuracy: Monotonicity with group size and, as  $N \rightarrow \infty$

$$p_{0|1}(N, \text{majority}) \rightarrow \begin{cases} 0, & \text{if } p_{0|1} < 1/2 \\ 1, & \text{if } p_{0|1} > 1/2 \\ \sqrt{N/(2\pi)} (4p_{0|1})^{\frac{N}{2}-1}, & \text{if } p_{0|1} < 1/4 \end{cases}$$

**Expected Decision Time:** Assume  $p_{1|1} > p_{0|1}$  and define

$$t_{<\frac{1}{2}} := \max\{t \in \mathbb{N} \mid p_{1|1}(0) + \dots + p_{1|1}(t) < 1/2\},$$

$$t_{>\frac{1}{2}} := \min\{t \in \mathbb{N} \mid p_{1|1}(0) + \dots + p_{1|1}(t) > 1/2\}$$

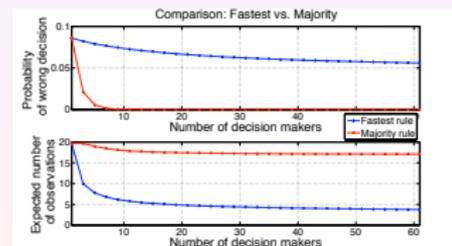
Then

$$\lim_{N \rightarrow \infty} \mathbb{E}[T | H_1, N, \text{majority}] = \frac{1}{2} (t_{<\frac{1}{2}} + t_{>\frac{1}{2}} + 1)$$

**Accuracy:** Monotonicity with group size and, as  $N \rightarrow \infty$

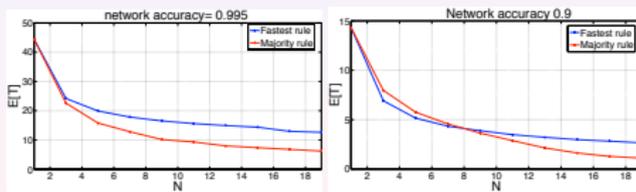
$$p_{0|1}(N, \text{majority}) \rightarrow \begin{cases} 0, & \text{if } p_{0|1} < 1/2 \\ 1, & \text{if } p_{0|1} > 1/2 \\ \sqrt{N/(2\pi)} (4p_{0|1})^{\frac{N}{2}}, & \text{if } p_{0|1} < 1/4 \end{cases}$$

	Accuracy	Expected decision time
<b>Fastest</b>	SDM accuracy at $t_{min}$	earliest possible decision time $t_{min}$
<b>Majority</b>	exponentially better than SDM	average of half-times $t_{<\frac{1}{2}}, t_{>\frac{1}{2}}$



## A fair comparison

- to compare different thresholds, re-scale local accuracy
- the group accuracy is now same (eg, low or high)
- compare the decision time



for most cases **majority** rule is best  
for some small inaccurate networks, **fastest** rule is best

## Conclusions and future directions



**Summary** fundamental understanding of “sequential aggregation”

- applicable to broad range of agent models, eg, mixed networks
- applicable to family of threshold-based rules
- tradeoffs in fastest vs majority
- role of time in sequential aggregation

**Future directions**

- models with heterogeneous agents
- models with interactions between agents
- models with correlated information
- how to use this analysis for design