

Boundary patrol using robotic networks without localization

Topology and Minimalism
RSS Workshop, June 28, 2008

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Incomplete state of the art



AeroVironment Inc, "Raven"
small unmanned aerial vehicle



iRobot Inc, "PackBot" unmanned
ground vehicle

Distributed algorithms

automata-theoretic: "Distributed Algorithms" by N. Lynch, D. Peleg
numerical: "Parallel and Distributed Computation" by Bertsekas and Tsitsiklis

Motion coordination

"rendezvous" by Suzuki and Yamashita
"consensus, flocking, agreement" by Jadbabaie, Olfati-Saber
"formation control" by Baillieul, Morse, Anderson

Research directions

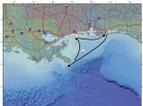
Build: distributed systems
embedded actuator/sensors networks

Develop distributed disciplines:

- (i) sensor fusion
- (ii) communications
- (iii) coordinated control
- (iv) task allocation and scheduling

Challenges

- (i) scalability
- (ii) performance
- (iii) robustness
- (iv) models



Environmental monitoring



Building monitoring and evac



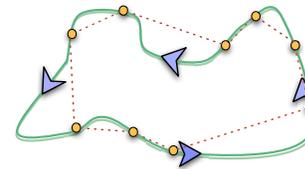
Security systems

Scenario 1: Boundary estimation

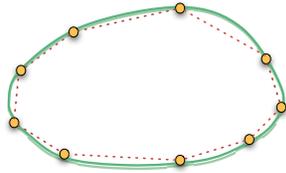
Assumption: local sensing and tracking, interpolation via waypoints

Objective: estimate/interpolate moving boundary

adaptive polygonal approximation



Scenario 1: Interpolation theory

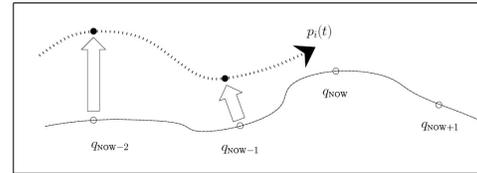


For strictly convex bodies (Gruber '80)

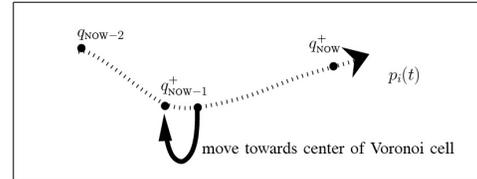
- **sufficient condition for optimality:** each two consecutive interpolation points p_k, p_{k+1} are separated by same line integral $\int_{p_k \rightarrow p_{k+1}} \kappa(\ell)^{1/3} d\ell$
- **error estimate** $\approx \frac{1}{12n^2} \left(\int_{\partial Q} \kappa(\ell)^{1/3} d\ell \right)^3$

Scenario 1: Estimate-Update and Pursuit

(i) projection step

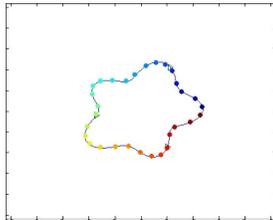


(ii) update interpolation points for "pseudo-uniform" interpolation placement



(iii) accelerate/decelerate for uniform vehicle placement

Scenario 1: Performance/robustness



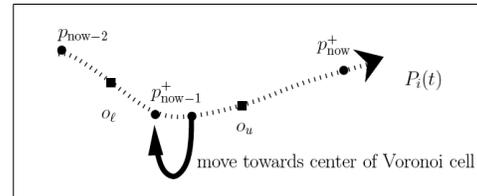
- asynchronous distributed over ring
- convergence to equally distributed interpolation points and equally spaced vehicles
- time complexity: worst case $O(n^2 \log(n))$, where $n = \frac{\# \text{ interpolation points}}{\# \text{ vehicles}}$
- ISS robust to: evolving boundary, interpolation, sensor noise

Scenario 1: translation into average consensus

- pseudo-distance between interpolation points (p_k, p_{k+1})

$$d(k) = \lambda \int_{p_k \rightarrow p_{k+1}} \kappa(\ell)^{1/3} d\ell + (1 - \lambda) \int_{p_k \rightarrow p_{k+1}} d\ell$$

- "go to center of Voronoi cell" update is peer-to-peer averaging rule

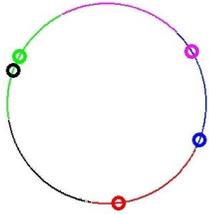


- linear model is:
 - stochastic matrices: switching, symmetric and nondegenerate
 - union of associated graphs over time is a ring (i.e., jointly connected graphs)
 - convergence rate as in Toeplitz tridagonal problem

Scenario 2: Synchronized boundary patrolling

- (i) some UAVs move along boundary of sensitive territory
- (ii) short-range communication and sensing
- (iii) surveillance objective:
 - minimize service time for appearing events
 - communication network connectivity

Example motion:



joint work with: Susca, Martínez

Analogy with mechanics and dynamics

- (i) robots with “communication impacts” analogous to **beads on a ring**
- (ii) classic subject in dynamical systems and geometric mechanics
 - billiards in polygons, iterated impact dynamics, gas theory of hard spheres
- (iii) rich dynamics with even just 3 beads (distinct masses, elastic collisions)
 - dynamics akin billiard flow inside acute triangle
 - dense periodic and nonperiodic modes, chaotic collision sequences

SIAM Review
Vol. 47, No. 2, pp. 273–300

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Iterated Impact Dynamics of N-Beads on a Ring*

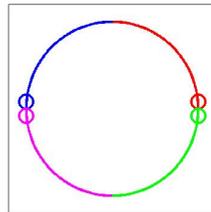
Bryan Cooley¹
Paul K. Newton¹

Abstract. When N -beads slide along a frictionless hoop, their collision sequence gives rise to a dynamical system that can be studied via matrix products. It is of general interest to understand the distribution of velocities and the corresponding eigenvalue spectrum that a given collision sequence can produce. We formulate the problem for general N and state some basic theorems regarding the eigenvalues of the collision matrices and their products. The

Boundary patrolling: synchronized bead oscillation

Desired synchronized behavior:

- starting from random initial positions and velocities
- every bead impacts its neighbor at the same point
- simultaneous impacts



Design specification for synchronization algorithm

Achieve: asymptotically stabilize synchronized motion

Subject to:

- (i) arbitrary initial positions, velocities and directions of motion
- (ii) beads can measure traveled distance, however
 - no absolute localization capability, no knowledge of circle length
- (iii) no knowledge about n , adaptation to changing n (even and odd)
- (iv) anonymous agents with memory and message sizes independent of n
- (v) smooth dependency upon effect of measurement and control noise

Fully-adaptive feedback synchronization



Algorithm: (for presentation's sake, beads sense their position)

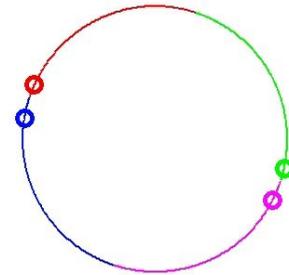
1st phase: compute **average speed** v and **desired sweeping arcs**

2nd phase for $f \in]\frac{1}{2}, 1[$, each bead:

- moves at nominal speed v if inside its desired sweeping arc
- slows down (fv) when moving away of its desired sweeping arc
hesitate when early
- when impact, change direction
- speeds up when moving towards its desired sweeping arc

Balanced initial condition:

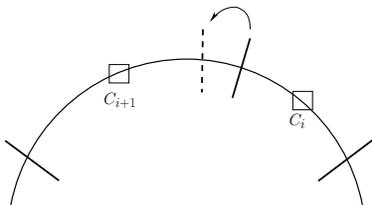
- n is even
- d_i is direction of motion
- $\sum_i^n d_i(0) = \sum_i^n d_i(t) = 0$
- $n/2$ move initially clockwise



First phase: average speed and sweeping arc

If an impact between bead i and $i + 1$ occurs:

- beads average nominal speeds: $v_i^+ = v_{i+1}^+ = 0.5(v_i + v_{i+1})$
- beads change their direction of motion if $d_i = -d_{i+1}$ (**head-head type**)
- beads update their desired sweeping arc



exponential average consensus

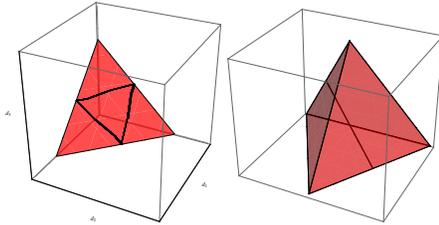
Challenges

- how to prove balanced synchronization?
- what happens for unbalanced initial conditions $\sum_i^n d_i(0) \neq 0$?
- what happens for n is odd?
- how to describe the system with a single variable?

Modeling detour

- configuration space

- (i) order-preserving dynamics $\theta_i \in \text{Arc}(\theta_{i-1}, \theta_{i+1})$ on \mathbb{T}^n
- (ii) $\Delta^n \times \{c, cc\}^n \times (\mathbb{R}_{>0})^n \times (\text{arcs})^n \times \{\text{away, towards}\}^n$

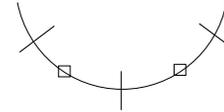


- hybrid system with

- (i) piecewise constant dynamics
- (ii) event-triggered jumps: impact, cross boundary

Passage and return times

- passage time:** $t_i^k = k$ th time when bead i passes by sweeping arc center



- return time:** $\delta_i(t) =$ duration between last two passage times
- if impact between beads $(i, i + 1)$ at time t , then

$$\begin{bmatrix} \delta_i \\ \delta_{i+1} \end{bmatrix} (t^+) = \underbrace{\begin{bmatrix} \frac{1-f}{1+f} & \frac{2f}{1+f} \\ \frac{2f}{1+f} & \frac{1-f}{1+f} \end{bmatrix}}_{\text{stochastic (irr + aperd)}} \begin{bmatrix} \delta_i \\ \delta_{i+1} \end{bmatrix} (t^-)$$

Convergence results: balanced synchronization

Balanced Synchronization Theorem: For balanced initial directions, assume

- (i) exact average speed and desired sweeping arcs
- (ii) initial conditions lead to well-defined 1st passage times

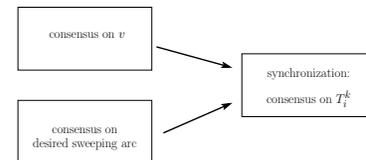
Then balanced synchronization is asymptotically stable

$$\lim_{t \rightarrow \infty} \delta(t) = \frac{2\pi}{Nv} \mathbf{1}_n, \quad \lim_{k \rightarrow +\infty} \left\| T^k - \frac{\mathbf{1}_n \cdot T^k}{n} \mathbf{1}_n \right\| = 0$$

Conjectures arising from simulation results

Only assumption required is balanced initial conditions.

- (i) analysis of cascade consensus algorithms

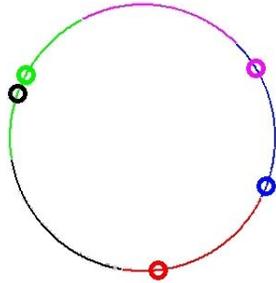


- (ii) global attractivity of synchronous behavior

Simulations results: 1-unbalanced case

1-unbalanced initial condition:

- n is odd
- $\sum_i^n d_i(0) = \sum_i^n d_i(t) = \pm 1$



1-unbalanced synchronization

- $f \in]\frac{1}{2}, \frac{n}{1+n}[$
- 1-unbalanced sync: beads meet at arcs boundaries $\pm \frac{2\pi}{n^2} \frac{f}{1-f}$

1-unbalanced Synchronization Theorem: For $\sum_i^n d_i(0) = \pm 1$, assume

- exact average speed and desired sweeping arcs
- initial conditions lead to well-defined 1st passage times

Then 1-unbalanced synchronization is asymptotically stable

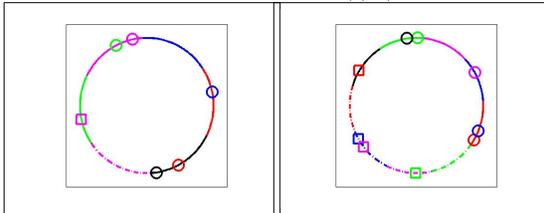
$$\lim_{t \rightarrow \infty} \delta(t) = \frac{2\pi}{Nv} \mathbf{1}_n, \quad \lim_{k \rightarrow +\infty} (T^{2k} - T^{2(k-1)}) = \frac{2}{v} \frac{2\pi}{n} \mathbf{1}_n$$

General unbalanced case

Conjecture global asy-synchronization in the balanced and unbalanced case

D -unbalanced period orbits Theorem:

Let $\sum_i^n d_i(0) = \pm D$. If there exists an orbit along which beads i and $i+1$ meet at boundary $\pm \frac{2\pi}{n^2} \frac{f}{1-f}$, then $f < \frac{n/|D|}{1+n/|D|}$.



Emerging discipline: motion-enabled networks

- **network modeling**
network, ctrl+comm algorithm, task, complexity
- **coordination algorithm**
deployment, task allocation, boundary estimation

Open problems

- algorithmic design for motion-enabled sensor networks
scalable, adaptive, asynchronous, agent arrival/departure tasks: search, exploration, identify and track
- integration between motion coordination, communication, and estimation tasks
- Very few results available on:
 - scalability analysis: time/energy/communication/control
 - robotic networks over random geometric graphs
 - complex sensing/actuation scenarios